An introduction to fractional calculus

Fundamental ideas and numerics



Fabio Durastante

Università di Pisa fabio.durastante@unipi.it fdurastante.github.io

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All-at-once

We have seen that for a problem of the form

$$egin{aligned} & u : \ \Omega imes [0, T] o \mathbb{R}^d, \ \Omega \subseteq \mathbb{R}^d \ & u(\mathbf{x}, 0) = u_0(\mathbf{x}), \ & \mathcal{B}(u) = 0, \qquad \mathbf{x} \in \partial \Omega. \end{aligned}$$

with

- $\mathcal{L}(\cdot)$ a *linear* and *autonomous* differential operator (possibly involving fractional derivatives),
- \rightarrow or changing u_t with ${}^{CA}D^{\alpha}_{[0,t]}u$,

we can rewrite it as a single linear system/matrix equation.

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$$egin{aligned} & \mathbf{u}_t = \mathcal{L}_h(\mathbf{u}), \quad \mathbf{u} : \ \mathbb{R}^n imes [0, \ \mathcal{T}]
ightarrow \mathbb{R}^n \ & \mathbf{u}(0) = \mathbf{u}_0, \ & \mathcal{B}_h(\mathbf{u}) = \mathbf{0}. \end{aligned}$$

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we can rewrite it as a single linear system/matrix equation.

To abstract the procedure let's think about working the Method Of Line way!

All-at-once: system of autonomous ODE

Following the MOL trail, we now have to solve a system of autonomous ODEs:

$$M\mathbf{u}_t(t) = L\mathbf{u}(t), \qquad M, L \in \mathbb{R}^{n \times n},$$

 \rightarrow that could be a differential-algebraic system of equations (DAE) if det(M) = 0.

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 To formulate the *all-at-once* procedure, one has to select a method to *march in time* the solution:

- 差 Linear multistep methods,
- 差 Runge-Kutta methods,
- ✗ General linear methods (a mix of the two above strategies).

Given a general ODE of the form

$$u'(t) = f(t, u(t)), \quad u(t_0) = u_0,$$

a k-step LMM is a recursion of the form with step-size $h = t_{n+k} - t_{n+k-1} > 0$

$$\sum_{j=0}^{k} \alpha_j u_{n+j} = \sum_{j=0}^{k} h \beta_j f_{n+j}, \qquad f_m \triangleq f(t_m, y_m),$$

with coefficients $\alpha_j \in \mathbb{R}$ and $\beta_j \in \mathbb{R}$ (j = 0, ..., k), and we are **interested only** in **implicit methods**, i.e., $\beta_k \neq 0$.

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They can be analyzed by looking at the polynomials

$$\rho(\zeta) = \sum_{j=0}^{k} \alpha_j \zeta^j = (\zeta - 1) \sum_{j=0}^{k-1} \gamma_j \zeta^j = (\zeta - 1) \cdot \rho_R(\zeta), \qquad \sigma(\zeta) = \sum_{j=0}^{k} \beta_j \zeta^j.$$

0-stable method

A method is 0-stable if all roots of $\rho(\zeta) = (\zeta - 1) \cdot \rho_R(\zeta) = 0$ lie inside or on the unit circle, with no multiple unimodular roots.

Zero stability is necessary for convergence,
 k It is a condition on the *extraneous operator* ρ_R(ζ), i.e., a condition on the k coefficients {γ_j}^{k-1}_{j=0}.

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A-stable method

The behavior of these methods can be analyzed by applying them on the test problem y' = ky subject to the initial condition y(0) = 1 with $k \in \mathbb{C}$. The solution of this equation is $y(t) = e^{kt}$. If the numerical method exhibits the same behavior of the solution for a fixed step size, then the method is said to be *A*-stable.

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B Usually one ends up with limitations involving the admissible h.

If we use a LMM with k > 1 we need more starting values than the one we have! We are interested in **diffusion dominated problems**, thus **B**ackward-**D**ifferentiation Formulas are a common choice.

$\{lpha_k\}_k$, $eta_k=1$, $eta_j=0$, $j\leq k$							
BDF2					1/2	-2	3/2
BDF3				-1/3	3/2	-3	11/6
BDF4			1/4	-4/3	3	—4	25/12
BDF5		-1/5	5/4	-10/3	5	-5	137/60
BDF6	1/6	-6/5	15/4	-20/3	15/2	-6	147/60



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$\{lpha_k\}_k,\ eta_k=1,\ eta_j=0,\ j\leq k$	4
BDF2 1/2 -2 3/2	2
BDF3 -1/3 3/2 -3 11/6	o
BDF4 1/4 -4/3 3 -4 25/12	-2
BDF5 -1/5 5/4 -10/3 5 -5 137/60	-4
BDF6 $1/6$ $-6/5$ $15/4$ $-20/3$ $15/2$ -6 $147/60$	

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Ne can use lower order BDFs to generate the step we need.

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BDF2 1/2 -2 3/2	2 4
BDF3 $-1/3 \ 3/2 \ -3 \ 11/6$	6 ²
BDF4 1/4 -4/3 3 -4 25/12	2 -2
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$\{lpha_k\}_k,\;eta_k=1,\;eta_j=0,j\leq k$	15
BDF2 1/2 -2 3/2	10
BDF3 -1/3 3/2 -3 11/6	0
BDF4 1/4 -4/3 3 -4 25/12	-5
BDF5 -1/5 5/4 -10/3 5 -5 137/60	-10
BDF6 $1/6$ $-6/5$ $15/4$ $-20/3$ $15/2$ -6 $147/60$	-15 10 -5 0 5 10 15 20 25

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$\{lpha_k\}_k,\;eta_k=1,\;eta_j=0,j\leq k$	30
BDF2 1/2 -2 3/2	20
BDF3 -1/3 3/2 -3 11/6	0
BDF4 1/4 -4/3 3 -4 25/12	-10
BDF5 -1/5 5/4 -10/3 5 -5 137/60	-20
BDF6 1/6 -6/5 15/4 -20/3 15/2 -6 147/60	$-30\frac{1}{20}$ -10 0 10 20 30 40

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From what we have seen in the last lectures we can write down the problem as

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$$(A_m \otimes M_n - h B_m \otimes L_n)\mathbf{u} = \mathbf{f},$$

$$B_m = \begin{bmatrix} 1 & & & \\ & 1 & & \\ & & \ddots & \\ & & & 1 \end{bmatrix}_{(m-1) \times (m-1)}$$

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$$(A_m \otimes M_n - hB_m \otimes L_n)\mathbf{u} = \mathbf{f},$$

$$\mathbf{f} = \begin{bmatrix} \mathbf{u}_0 + f(t_1) \\ -1/2\mathbf{u}_0 + f(t_2) \\ 1/3\mathbf{u}_0 + f(t_3) \\ -1/4\mathbf{u}_0 + f(t_4) \\ 1/5\mathbf{u}_0 + f(t_5) \\ -1/6\mathbf{u}_0 + f(t_5) \\ -1/6\mathbf{u}_0 + f(t_6) \\ f(t_7) \\ \vdots \end{bmatrix}$$

```
v0 = [1:1]:
n = length(L);
% Discretize
m = 100:
T = linspace(0, 10, m); h = T(2) - T(1);
r = zeros(m-1,1); c = zeros(m-1,1);
r(1:7) = [147/60, -6, 15/2, -20/3, 15/4, -6/5, 1/6]:
c(1) = 147/60:
A = toeplitz(r,c);
A(1,1) = 1; % Fix BCs
A(2,1) = -2; A(2,2) = 3/2;
A(3,1) = 3/2; A(3,2) = -3; A(3,3) = 11/6;
A(4,1) = -4/3; A(4,2) = 3; A(4,3) = -4;
\rightarrow A(4,4) = 25/12;
A(5,1) = 5/4; A(5,2) = -10/3; A(5,3) = 5;
```

```
A(5,4) = -5; A(5,5) = 137/60;
In = speye(n,n);
Im = speye(m-1, m-1);
%% Build rhs:
b = zeros((m-1)*n, 1):
b(1:2) = v0;
b(3:4) = -1/2*y0;
b(5:6) = 1/3*y0;
b(7:8) = -1/4*y0;
b(9:10) = 1/5*v0;
b(11:12) = -1/6*y0;
% SOLVE (Linear system)
M = kron(A, In) - h * kron(Im, L);
x = M \setminus b:
```

We can compare the solution with ode15s, and visualize it

```
[tt,yy] = ode15s(@(t,y) L*y,T,y0);
X = reshape(x,n,m-1);
X = [y0,X];
% Plot
plot(T,X(1,:),'r-',T,X(2,:),'b-',...
T,yy(:,1),'ro',...
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 A_m is a banded Toeplitz matrix plus a rank correction.

$$A_{m} = \begin{bmatrix} 1 & & & & \\ -2 & 3/2 & & & \\ 3/2 & -3 & 11/6 & & & \\ -4/3 & 3 & -4 & 25/12 & & \\ 5/4 & -10/3 & 5 & -5 & 137/60 & & \\ -6/5 & 15/4 & -20/3 & 15/2 & -6 & 147/60 & \\ 1/6 & -6/5 & 15/4 & -20/3 & 15/2 & -6 & 147/60 & \\ \vdots & \ddots \end{bmatrix}$$

Solution We know the eigenvalues in closed form: it's lower triangular!

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- We know the eigenvalues in closed form: it's lower triangular!
- The Field-Of-Values contains the origin... bad for bounds!
- 📅 Its clearly non diagonalizable, if we try and look at the condition number of the eigenvector matrix $\kappa_2(X_{100}) = 7.30 \times 10^{111}$.

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$A_m =$	$\begin{bmatrix} 1\\ -1 \end{bmatrix}$	1 · · -	·. -1	1 $(m-1) \times (m-1)$
		BDF	1, 0	$\kappa = 1.5$
	т	n	IT	Residual
	64	128	13	1.007848e-10
	128	256	16	6.145733e-10
	256	512	21	7.639171e-10
	512	1024	27	5.857467e-10
	1024	2048	34	8.065585e-10
	2048	4096	42	9.819085e-10

It is a Jordan block, so *no diagonalization*,
What do we expect for the matrix equation solver?

BDF6, $lpha=1.5$					
т	п	IT	Residual		
64	128	21	3.651570e-10		
128	256	33	1.746513e-10		
256	512	71	2.530720e-15		
512	1024	128	1.975160e-22		
1024	2048	251	4.157259e-10		
2048	4096	495	6.310887e-10		

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- If we use more than one step, we still need auxiliary formulas to close the iteration.
- We could distribute the conditions differently, that is, not all on the initial data.

$$\sum_{j=-\nu}^{\mu-\nu} \alpha_{j+\nu} \mathbf{u}_{n+j} = h \sum_{j=-\nu}^{\mu-\nu} \beta_{j+\nu} \mathbf{f}_{n+j}, \quad n = \nu, \dots, m-k+\nu.$$

🔗 k steps,

- ${\ensuremath{ @ } \ensuremath{ \partial \mu } \nu}$ final conditions,

$$\textit{ O} Escribed by \ \rho(z) = z^{\nu} \sum_{j=-\nu}^{k-\nu} \alpha_{j+\nu} z^{j}, \text{ and } \sigma(z) = z^{\nu} \sum_{j=-\nu}^{k-\nu} \beta_{j+\nu} z^{j}.$$

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How does this change matrices and stability?

If we collect the matrices for the inner steps of a scalar ODE, we get

$$A_{m} = \begin{bmatrix} \alpha_{\nu} & \cdots & \alpha_{k} & & & \\ \vdots & \ddots & & \ddots & & \\ \alpha_{0} & & \ddots & & \ddots & \\ & \ddots & & \ddots & & \alpha_{k} \\ & & \ddots & & \ddots & \vdots \\ & & & & \alpha_{0} & \cdots & \alpha_{\nu} \end{bmatrix}, B_{m} = \begin{bmatrix} \beta_{\nu} & \cdots & \beta_{k} & & & \\ \vdots & \ddots & & \ddots & & \\ \beta_{0} & & \ddots & & \ddots & \\ & \ddots & & \ddots & & \beta_{k} \\ & & \ddots & & \ddots & & \beta_{k} \\ & & & \ddots & & \ddots & \vdots \\ & & & & & \beta_{0} & \cdots & \beta_{\nu} \end{bmatrix}_{(m-\nu) \times (m-\nu)}$$

and the vectors

$$\mathbf{u} = (u_{\mathbf{v}}, \cdots, u_{m-1})^T, \quad \mathbf{f} = (f_{\mathbf{v}}, \cdots, f_{m-1})^T.$$

If we collect the matrices for the inner steps of a *scalar* ODE, we get A_m , B_m , and the vectors

$$\mathbf{u} = (u_{\mathbf{v}}, \cdots, u_{m-1})^T, \quad \mathbf{f} = (f_{\mathbf{v}}, \cdots, f_{m-1})^T.$$

Finding the system

 $A_{m}\mathbf{u}$ –

$$hB_m \mathbf{f} = -\begin{bmatrix} \sum_{j=0}^{\nu-1} (\alpha_j y_j - h\beta_j f_j) & \vdots \\ \vdots & \\ a_0 y_{\nu-1} - h\beta_0 f_{\nu-1} \\ 0 \\ \vdots \\ 0 \\ \alpha_k y_m - h\beta_k f_m \\ \vdots \\ \sum_{j=1}^{\mu} (\alpha_{\nu+j} y_{m-1+j} - h\beta_{\nu_1+j} f_{m-1+j}). \end{bmatrix}$$

- A_m and B_m are Toeplitz matrices with *lower* bandwidth ν and upper bandwidth μ.
- We still need auxiliary formulas to fix the $v + \mu - 1$ starting/ending values.

Before concluding the construction, let's focus on *convergence* and *stability*.

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 $S_{\nu,\mu}$ -polynomial (Brugnano and Trigiante 1998, Definition 4.4.2)

A polynomial p(z) of degree $k = v + \mu$ is an $S_{v,\mu}$ -polynomial if its roots are such that

 $|z_1| \leq |z_2| \leq \cdots \leq |z_\nu| < 1 < |z_{\nu+1}| \leq \cdots \leq |z_\nu|.$

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 \odot If $\nu = k$ ($\mu = 0$), these are the conditions for LMF 0-stability!

Let $a_{-\nu}a_{\mu} \neq 0$ and



we consider the polynomial

$$p(z) = \sum_{i=-\nu}^{\mu} a_i z^{\nu+i}.$$

Let $a_{-\nu}a_{\mu} \neq 0$ and

$$T_n = \begin{bmatrix} a_0 & \cdots & a_{\mu} & & & \\ \vdots & \ddots & & \ddots & & \\ a_{-\nu} & & \ddots & & \ddots & & \\ & \ddots & & \ddots & & \ddots & \\ & & \ddots & & \ddots & & a_{\mu} \\ & & & \ddots & & \ddots & \vdots \\ & & & & a_{-\nu} & \cdots & a_0 \end{bmatrix},$$

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Lemma (Brugnano and Trigiante 1998, Lemma 4.4.4)

If the polynomial p(z) associated with the matrix T_n is an $N_{\nu,\mu}$ -polynomial, then T_n^{-1} has entries $t_{i,j}^{(-1)}$ such that 1. $|t_{i,j}^{(-1)}| \leq \gamma$ independent of N, for $i \geq j$, 2. $|t_{i,j}^{(-1)}| \leq \eta \xi^{j-i}$ for i < j, where $\eta > 0$ and $0 < \xi < 1$ are independent of N.

Let $a_{-\nu}a_{\mu} \neq 0$ and

$$T_{n} = \begin{bmatrix} a_{0} & \cdots & a_{\mu} & & \\ \vdots & \ddots & & \ddots & \\ a_{-\nu} & & \ddots & & \ddots & \\ & \ddots & & \ddots & & a_{\mu} \\ & & \ddots & & \ddots & & a_{\mu} \\ & & & \ddots & & \ddots & \vdots \\ & & & & a_{-\nu} & \cdots & a_{0} \end{bmatrix},$$

we consider the polynomial

$$p(z) = \sum_{i=-\nu}^{\mu} a_i z^{\nu+i}.$$

Lemma (Brugnano and Trigiante 1998, Lemma 4.4.4)

If the polynomial p(z) associated with the matrix T_n is an $N_{\nu,\mu}$ -polynomial, then T_n^{-1} has entries $t_{i,j}^{(-1)}$ such that

 |t_{i,j}⁽⁻¹⁾| ≤ γ independent of N, for i ≥ j,
 |t_{i,j}⁽⁻¹⁾| ≤ ηξ^{j-i} for i < j, where η > 0 and 0 < ξ < 1 are independent of N.

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Let $a_{-\nu}a_{\mu} \neq 0$ and

$$T_n = egin{bmatrix} a_0 & \cdots & a_\mu & & & \ dots & \ddots & & \ddots & & \ a_{-\nu} & & \ddots & & \ddots & & \ & \ddots & & \ddots & & a_\mu \ & & \ddots & & \ddots & & a_\mu \ & & & \ddots & & \ddots & dots & \ & & & a_{-\nu} & \cdots & a_0 \end{bmatrix},$$

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and $0 < \xi < 1$ are independent of *N*.

with Δ_n the upper triangular Toeplitz matrix with last column $(\xi^{n-1}, \ldots, \xi^2, \xi, 0)^T$.

Theorem (Brugnano and Trigiante 1998, Theorem 4.4.3)

Ignoring the effect of round-off errors, a BVM with (ν, μ) -boundary conditions is convergent if it is consistent and the polynomial $\rho(z)$ is an $N_{\nu,\mu}$ -polynomial.

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To reproduce the "0-stable + consistent \Rightarrow convergence" framework, we define:

$0_{\nu,\mu}$ -stability (Brugnano and Trigiante 1998, Definition 4.5.1)

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(ν, μ) -Absolute stability (Brugnano and Trigiante 1998, Definition 4.7.1)

A BVM with (ν, μ) -boundary conditions is ν, μ -Absolutely stable for a given complex number q it the polynomial $\pi(z, q) = \rho(z) - q\sigma(z)$, is an $S_{\nu,\mu}$ -polynomial.

We have a degree of **arbitrariness** in deciding how and how many initial / final conditions to set. Clearly ν has to be at least one (we do have an initial condition of our IVP), then for the remaining we have to let (ν, μ) -Absolute stability guide us.

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Correct use a consistent LMF is *correctly used* in $q \in \mathbb{C}^-$, where $\pi(z, q)$ is an $S_{\nu,\mu}$ -polynomial, if ν conditions are imposed at the initial points, and μ conditions are posed at the end of the interval.

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To have a livable life, one always consider family of methods for which the boundary of the (ν, μ) -Absolutely stability region is a *regular Jordan curve*. More specifically, having that

 $\mathcal{A}_{\nu,\mu} = \{ q \in \mathbb{C} : \pi(z,q) \text{ is an } S_{\nu,\mu}\text{-polynomial} \},$

has the origin on its boundary and is possibly equal to the whole \mathbb{C}^- .

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‡ BDF \Rightarrow Generalized-BDF (GBDF): $\sum_{i=0}^{k} \alpha_i u_{n+i} = hf_{n+j}, j \in \{0, 1, \dots, k\}$

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- **Proof** BDF \Rightarrow Generalized-BDF (GBDF): $\sum_{i=0}^{k} \alpha_i u_{n+i} = hf_{n+j}, j \in \{0, 1, \dots, k\}$
 - **()** A method of this form is $0_{\nu,k-\nu}$ -stable and $A_{\nu,k-\nu}$ -stable for

$$u = egin{cases} rac{k+2}{2}, & ext{for even } k, \ rac{k+1}{2}, & ext{for odd } k. \end{cases}$$

 \Rightarrow with this choice we no longer have the constraint of having at most k = 6 steps of the standard BDF!

‡ Adams-Moulton Methods \Rightarrow GAMM $u_{n+j} - u_{n+j-1} = h \sum_{i=0}^{k} \beta_i f_{n+i}$

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See the book (Brugnano and Trigiante 1998) for other possible generalizations.

→ We need additional formulas for the $k - 1 = v + \mu - 1$ boundary values.

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✓ We need additional formulas for the k − 1 = v + µ − 1 boundary values.
If we know how to compute them, then we end up having to solve the matrix equation

$$M_n U A_m^T - h L_n U B_m^T = F,$$

or the linear system

 $(A_m \otimes M_n - hB_m \otimes L_n)\mathbf{u} = \mathbf{f}$, where $\operatorname{vec}(U) = \mathbf{u}$, $\operatorname{vec}(F) = \mathbf{f}$.

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, where $\operatorname{vec}(U) = \mathbf{u}$, $\operatorname{vec}(F) = \mathbf{f}$.

Let us build everything for using GBDFs and our fractional-in-space problem.

First we need to compute $\rho(z)$ and $\sigma(z)$

```
function [ro,si] = rosi_bdf( k, j )
b = zeros(k+1,1); b(2) = 1;
ro = vsolve( -j:k-j, b(:) );
si = zeros( k+1, 1 ); si( j+1 ) = 1;
end
```

Coefficients are computed by imposing consistency of maximal order p:

$$\sum_{j=0}^k (j^s \alpha_j - s j^{s-1} \beta_j) = 0,$$

 $s=0,1,\ldots,p$.

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$$\sum_{j=0}^{k} (j^{s} \alpha_{j} - s j^{s-1} \beta_{j}) = 0,$$

$$s=0,1,\ldots,p$$
.

```
function f = vsolve(x, b)
f = b;
n = length(x) - 1;
for k = 1:n
 for i = n+1:-1:k+1
 f(i) = f(i) - x(k) * f(i-1):
 end
end
for k = n:-1:1
 for i = k+1 \cdot n+1
  f(i) = f(i)/(x(i) - x(i-k)):
 end
 for i = k:n
 f(i) = f(i) - f(i+1);
 end
end
end
```

Then we use the ro_si routine to build the A_m and B_m matrices

```
function [a,b] = mab(k, n)
nu = fix((k+2)/2);
a = spalloc(n, n+1, (k+1)*n);
b = a:
for i = 1:nu
 [ro,si] = rosi_bdf( k, i );
 a(i,1:k+1) = ro.':
 b(i.1:k+1) = si.':
end
for i = nu+1:n-(k-nu)
 a(i,i+1+(-nu:k-nu)) = ro.':
 b(i,i+1+(-nu:k-nu)) = si.':
end
```

```
j = nu;
for i = n-(k-nu)+1:n
 i = i + 1;
 [ro,si] = rosi_bdf( k, j );
 a(i.n+1+(-k:0)) = ro.':
 b(i,n+1+(-k:0)) = si.';
end
end
\Rightarrow for i = 1:nu: end. initial conditions.
\langle \rangle for i = nu+1:n-(k-nu); end.
    Toepltiz part,
\downarrow for i = n-(k-nu)+1:n; end, final
    conditions
```

We can use the routine to generate [Alpha,Beta] = mab(k,m); A = Alpha(:,2:m+1); B = Beta(:,2:m+1); and visualize them



• The first column contains the coefficients needed to compute the right-hand-side.

We now need to build the right-hand-side

```
nk=n*(m+1);
b=zeros(nk,1); % Allocate the space for one more than needed
for j=1:m % Use the source to build the rhs:
    b(1+j*n:(j+1)*n)=f(x,t0+j*h);
end
b(n+1:n*(m+1))=h*kron(Beta,speye(n))*b; % Correct with the betas coeff.s
b(1:n)=u0; % First block as the initial condition
% Correction coefficients:
Am = kron(Alpha(:,1),speye(n))-h*kron(Beta(:,1),L);
b(n+1:nk)=b(n+1:nk)-Am*u0; % Finish building RHS
```

And then we can solve the linear system

```
Mat = kron(A,M) - h*kron(B,L); rhs = b(n+1:nk);
u = Mat\rhs;
```

We can compare the solution with ode15s:

```
U = [u0,reshape(u,n,m)]; t = t0:h:tf;
[TT,UU] = ode15s(@(t,y) L*y +

→ f(x.',t),t,u0);
E = abs(U-reshape(UU,m+1,n).');
figure(2)
subplot(1,3,1)
mesh(t,x,U);
xlabel('t');
ylabel('t');
title('GBDF(6,100) on 100')
```

```
subplot(1,3,2)
mesh(t,x,reshape(UU,m+1,n).')
xlabel('t');
ylabel('x');
title('ode15s')
subplot(1,3,3)
mesh(t,x,log10())
xlabel('t');
ylabel('x');
title('Error')
```

GBDF(6,100) on 100 ode15s Error -2 6 -3 4 -4 -5 2 2 -6 0 2 02 2 0.5 0.5 0.5 0 0 0 0 0 х х х

We can compare the solution with ode15s:

What happens if we attempt solution via our matrix-equation solver?

We can solve it by doing:

```
maxit = 100;
tol = 1e-9;
[LL,UL] = lu(-h*L);
[LA,UA] = lu(A);
[X1,X2,res]=kpik_sylv(-h*L,LL,UL,A,
\[code] LA,UA,C1,C2,maxit,tol);
```

Using our non-symmetric test problem with variable coefficients and fractional order α .
We can solve it by doing:

maxit = 100; tol = 1e-9; [LL,UL] = lu(-h*L); [LA,UA] = lu(A); [X1,X2,res]=kpik_sylv(-h*L,LL,UL,A, \[arrow LA,UA,C1,C2,maxit,tol);

k	т	п	IT	Res.
2	32	64	16	1.08e-15
2	64	128	23	2.16e-10
2	128	256	30	4.72e-10
2	256	512	38	9.20e-10
2	512	1024	49	7.31e-10
2	1024	2048	62	7.82e-10
2	2048	4096	78	8.06e-10
2	4096	8192	97	9.24e-10

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k	т	п	IT	Res.
3	32	64	15	7.18e-10
3	64	128	20	9.80e-10
3	128	256	26	7.77e-10
3	256	512	34	4.21e-10
3	512	1024	43	5.75e-10
3	1024	2048	54	8.05e-10
3	2048	4096	68	8.84e-10
3	4096	8192	85	9.87e-10

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k	т	п	IT	Res.
4	32	64	16	1.19e-14
4	64	128	24	3.22e-10
4	128	256	31	4.05e-10
4	256	512	39	6.97e-10
4	512	1024	50	6.20e-10
4	1024	2048	63	7.70e-10
4	2048	4096	79	9.05e-10
4	4096	8192	99	9.05e-10

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k	т	п	IT	Res.
5	32	64	16	1.72e-14
5	64	128	22	2.96e-10
5	128	256	28	4.90e-10
5	256	512	36	5.56e-10
5	512	1024	46	5.53e-10
5	1024	2048	58	7.10e-10
5	2048	4096	73	8.04e-10
5	4096	8192	91	9.75e-10

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k	т	п	IT	Res.
6	32	64	16	3.46e-14
6	64	128	24	4.70e-10
6	128	256	31	5.73e-10
6	256	512	40	4.78e-10
6	512	1024	50	9.39e-10
6	1024	2048	64	7.69e-10
6	2048	4096	81	7.31e-10
6	4096	8192	100	1.10e-09

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k	т	п	IT	Res.
7	32	64	16	6.13e-15
7	64	128	22	6.60e-10
7	128	256	29	4.78e-10
7	256	512	37	7.04e-10
7	512	1024	47	8.47e-10
7	1024	2048	60	7.66e-10
7	2048	4096	76	7.36e-10
7	4096	8192	95	8.46e-10

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k	т	п	IT	Res.
8	32	64	16	2.46e-14
8	64	128	24	5.41e-10
8	128	256	31	7.57e-10
8	256	512	40	6.53e-10
8	512	1024	51	7.34e-10
8	1024	2048	65	6.98e-10
8	2048	4096	82	7.42e-10
8	4096	8192	100	1.56e-09

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• The solution seems to be robust with respect to k,

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Using our non-symmetric test problem with variable coefficients and fractional order α .

• The solution seems to be robust with respect to k,

• We still have a small increase with n and m.

k	т	п	IT	Res.
8	32	64	16	2.46e-14
8	64	128	24	5.41e-10
8	128	256	31	7.57e-10
8	256	512	40	6.53e-10
8	512	1024	51	7.34e-10
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Let's now look for a different approach.

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We can do matrix vector products with the system matrix without assembling the matrix:

```
function [y] = Mprod(A,B,L,h,x)
[sp1,~] = size(A);
[m,~] = size(L);
X = reshape(x,m,sp1);
Y = X*A' - h*(L*X*B');
y = reshape(Y,m*sp1,1);
end
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end
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The linear system is not symmetric: we can use either GMRES or Flexible-GMRES to solve it.

We just need to figure out a preconditioner.

The *P*idea is *again* using a preconditioner that has the same structure:

$$P=\breve{A}_m\otimes M_n-h\breve{B}_m\otimes \tilde{L}_n,$$

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 - A_m, B_m are Toeplitz + low-rank ⇒ Circulant or Fast-Transform preconditioners,
 L̃_n has the *quasi-Toeplitz structure* we have seen, so we can use some of the techniques we had already seen for this; (Bertaccini and Durastante 2018).

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- **?** How do we select the approximations \breve{A}_m , \breve{B}_m and \tilde{L}_n ?
 - A_m , B_m are Toeplitz + low-rank \Rightarrow Circulant or Fast-Transform preconditioners,
 - L_n has the *quasi-Toeplitz structure* we have seen, so we can use some of the techniques we had already seen for this; (Bertaccini and Durastante 2018).
- **>** It would be good to also have a **parallel way of applying the preconditioner**.

• If \breve{A}_m and \breve{B}_m are *circulant-like approximations* of the Toeplitz (+ "low rank") matrices A_m and B_m , and the mass matrix is the identity, then we can express the **eigenvalues** of P as

$$\phi_i - h \psi_i \lambda_j, \qquad i = 1, \dots, m, \quad j = 1, \dots, n,$$

where

 $\not F \{ \phi_i \}$ are the eigenvalues of the circulant–like approximation \breve{A} , $\not F \{ \psi_i \}$ are the eigenvalues of the circulant–like approximation \breve{B} , $\not F \{ \lambda_j \}$ are the eigenvalues of the selected approximation of J_n .

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where

 $\begin{array}{l} \checkmark \{ \varphi_i \} \text{ are the eigenvalues of the circulant-like approximation } \breve{A}, \\ \checkmark \{ \psi_i \} \text{ are the eigenvalues of the circulant-like approximation } \breve{B}, \\ \checkmark \{ \lambda_j \} \text{ are the eigenvalues of the selected approximation of } J_n. \end{array}$

What circulant-like approximation do we want?

? If \breve{A}_m and \breve{B}_m are *circulant-like approximations* of the Toeplitz (+ "low rank") matrices A_m and B_m , and the mass matrix is the identity, then we can express the **eigenvalues** of P as

$$\phi_i - h \psi_i \lambda_j, \qquad i = 1, \dots, m, \quad j = 1, \dots, n,$$

where

{φ_i} are the eigenvalues of the circulant–like approximation Ă,
 {ψ_i} are the eigenvalues of the circulant–like approximation B,
 {λ_j} are the eigenvalues of the selected approximation of J_n.
 What circulant-like approximation do we want?

An idea could be using Strang approximation (Gu et al. 2015)

$$P_{\mathfrak{s}} = \mathfrak{s}(A_m) \otimes I_m - h\mathfrak{s}(B_m) \otimes L_n,$$







- $\mathfrak{s}(B) \text{ can be built}$ analogously.
- \$\vec{s}(A)\$ is singular due to the consistency condition.
- It is a single 0 eigenvalue, so we can move it by a rank 1 perturbation: \$\vec{s}(.)\$.



What can we say about the clustering properties of this preconditioner?

Theorem (Bertaccini 2000, Theorem 4.1)

Let $\mathcal{M} = A_m \otimes I_n - hB_m \otimes L_n$ for an $A_{\nu,k-\nu}$ -stable formulae with k steps. Let P be the block circulant preconditioner

$$P = \breve{A}_m \otimes M_n - h\breve{B}_m \otimes L_n.$$

Then, for fixed $\delta > 0$, there exists $C_{\delta} \ge 0$, $m_{\delta} \ge k$ such that, for all $m \ge m_{\delta}$ (m+1 is the size of A and B),

$$P^{-1}M = I + M^{(1)}_{\delta} + M^{(2)}_{\delta}$$

where $\operatorname{rank}(M_{\delta}^{(2)}) \leq n[2(k+1) + C_{\delta}]$ and $\|M_{\delta}^{(1)}\|_{2} \leq \delta c_{L}$ does not depend on *m*. If *P* is defined as Strang's circulant preconditioner, then $C_{\delta} = \|M_{\delta}^{(1)}\| = 0$.

Another available choice is using instead $\{\omega\}$ -Circulant matrices, i.e.,

$$P_{\omega} = \omega(A_m) \otimes I_n - h\omega(B_m) \otimes L_n,$$



- (a) $\omega(B_m)$ is defined similarly.
 - The usual choice is setting ω = -1, i.e., the skew-circulant preconditioner.

Stuctured preconditioner: application

To apply

$$P_{\omega}^{-1}\mathbf{v} = (\omega(A_m) \otimes I_n - h\omega(B_m) \otimes L_n)^{-1}\mathbf{v},$$

We can use the **diagonalization** of $\omega(A_m)$ and $\omega(B_m)$, i.e.,

$$P_{\omega}^{-1}\mathbf{v} = (F\Omega \otimes I_n)^{-1}(\Lambda_A \otimes I_n - h\Lambda_B \otimes L_n)^{-1}(\Omega^H F^H \otimes I_n)^{-1}\mathbf{v}.$$

1. Compute
$$\mathbf{w} = (\Omega^* F^* \otimes I_m)^{-1} \mathbf{v} = -V \Omega^{-H} F$$
,
2. Solve $(\Lambda_A \otimes I_n - h \Lambda_B \otimes L_n)^{-1} \mathbf{w}$ by solving

$$(\lambda_i(A)I_n - h\lambda_i(B)L_n)\mathbf{z}_i = \mathbf{w}_i, \quad i = 1, \dots, m$$

with $vec([\mathbf{w}_1,\ldots,\mathbf{w}_m]) = \mathbf{w}$, and similarly for \mathbf{z} ,

3. Compute $\mathbf{y} = (F\Omega \otimes I_n)^{-1}\mathbf{z} = -ZF^H\Omega^{-1}$.

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- 3. Compute $\mathbf{y} = (F\Omega \otimes I_n)^{-1}\mathbf{z} = -ZF^H\Omega^{-1}$.
- This step is embarrassingly parallel!

Numerical example

We use our favorite test problem with the space variant, nonsymmetric fractional operator in space and $\alpha = 1.5$, using GMRES(20) with a tolerance of 1e-9 using the P_{-1} preconditioner.

k = 2		k = 3		k = 4			k = 5			-	<i>k</i> = 6					
n	m	lt		n	m	lt	n	m	lt	n	m	lt	-	n	m	lt
64	32	30		64	32	32	64	32	35	64	32	38	-	64	32	46
128	64	31		128	64	33	128	64	38	128	64	45		128	64	53
256	128	31		256	128	34	256	128	39	256	128	48		256	128	58
512	256	31		512	256	34	512	256	39	512	256	50		512	256	62
1024	512	30		1024	512	33	1024	512	37	1024	512	49		1024	512	60

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• Reduced • iteration dependence, but paid with • full memory price!

Further modifications

We can further approximate the preconditioner by selecting instead of L_n in

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a suitable approximation, e.g.,

- $g_k(L_n)$ a bandwidth k approximation of the dense L_n matrix, i.e., using the information on the decay of the coefficients (Bertaccini and Durastante 2018).
- Structured preconditioner based on GLT theory.

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in Open areas of research

- **\dot{\mathbf{x}}** Efficient solution strategies for the $\lambda_i(A)I_n h\lambda_i(B)L_n$ systems,
- 🖈 Load-balancing issues for parallelism,
- * Optimal poles selection for the matrix-equation based solvers,
- **\dot{\mathbf{x}}** Multigrid solvers/preconditioners for $(A_m \otimes M_n hB_m \otimes L_n)\mathbf{u} = \mathbf{f}$.

‡ Tensor Equations

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\bigcirc Let us suppose that L_n is obtained as the discretization of a *multidimensional fractional operator*, i.e.,

$$L_n = \sum_{i=1}^{\ell} \left(\mathcal{K}_{m,\ell}^{-} \bigotimes_{p=1}^{i-1} I \otimes \mathcal{G}_{n^{1/\ell}}^{(\ell)} \otimes \bigotimes_{p=1}^{\ell-1} I + \mathcal{K}_{n,\ell}^{+} \bigotimes_{p=1}^{i-1} I \otimes \mathcal{G}_{n^{1/\ell}}^{(\ell)} {}^{\mathcal{T}} \otimes \bigotimes_{p=1}^{\ell-1} I \right)$$

where $K_{m,\ell}^{\pm}$ have also a Kronecker tensor structure whenever the functions $\{\kappa_j\}_{j=1}^{\ell}$ are separable in the x_i variables.

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where $K_{m,\ell}^{\pm}$ have also a Kronecker tensor structure whenever the functions $\{\kappa_j\}_{j=1}^{\ell}$ are separable in the x_j variables.

The matrix: $\mathcal{M} = A_m \otimes I_n - hB_m \otimes L_n$ has now a lot of redundant information!

As we have done for the hierarchical formats, we want

- \clubsuit A compressed representation of \mathcal{M} , possibly with a number of parameters that grows poly-logarithmically with the overall size...
- A fast BLAS-like toolbox to solve our problem in this format.
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There exists *many formats* for which this is possible, *e.g.*, the CANDECOMP/PARAFAC (CP) decomposition, the Tucker format, the Tensor Train (TT), the TT-Tucker, *etc.*; see (Kolda and Bader 2009).

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We focus on the \Box Tensor-Train format, since it has a simple enough toolbox to work with: \bigcirc TT-Toolbox.



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A tensor is a multidimensional array, $a \in \mathbb{R}$ is a 0-tensor, $\mathbf{v} \in \mathbb{R}^{n_1}$ is a 1-tensor, $A \in \mathbb{R}^{n_1 \times n_2}$ is a 2-tensor, $\mathcal{A} \in \mathbb{R}^{n_1 \times n_2 \times n_3}$ is a 3-tensor, ...

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- **E** A **tensor** is a **multilinear maps** with respect to a fixed finite-dimensional \mathbb{R} vector space V

$$\mathcal{A}: \underbrace{V^* imes \cdots imes V^*}_p imes \underbrace{V imes \cdots imes V}_q
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The definition we select depends on the operations we want to perform.

Tensor-Train¹

Let us start from trying to describe a *vector* associated with our discretization matrix \mathcal{M} .

¹For part of this material, a sincere thanks to Stefano Massei.

D Tensor-Train¹

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$$A(i_1, i_2) = U_1(i_1, :) \cdot U_2(:, i_2), \qquad A =$$

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where the two indices select the left and right vectors. In a tensor of order d we insert d-2 matrices between the two vectors:

$$\mathcal{T}(i_1,\ldots,i_d) = U_1(i_1, :) \cdot U_2(:, :, i_2) \cdot \ldots \cdot U_{d-1}(:, :, i_{d-1}) \cdot U_d(:, i_d)$$

$$\overbrace{k_j \atop n_i}^{k_j} \mathcal{T}(i_1,\ldots,i_d) = \overbrace{}$$

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More formally, a tensor ${\mathcal T}$ is in TT decomposition if it can be written as



- Smallest possible tuple (k_1, \ldots, k_{d-1}) is called the **TT-rank** of \mathcal{T} .
- $U_j \in \mathbb{C}^{k_{j-1} \times n_j \times k_j}$ are called the **TT** cores of \mathcal{T} (with $k_0 = k_d = 1$).
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If for any $1 \leq \mu \leq d-1$ we group the first μ factors and last $d-\mu$ factors then

 $\mathcal{T}(i_1,\ldots,i_{\mu},i_{\mu+1},\ldots,i_d),$

is the matrix-matrix product of two (large) matrices.

TT decomposition and matrix factorizations

The μ th unfolding of $\mathcal{T} \in C^{n_1 \times \cdots \times n_d}$ is obtained by arranging the entries in a matrix

 $\mathcal{T}^{<\mu>} \in \mathbb{C}^{(n_1 \cdots n_\mu) \times (n_{\mu+1} \cdots n_d)}$

where the corresponding index map is given by

$$\begin{array}{lll} \operatorname{ind} : \mathbb{N}^{n_1 \times \cdots \times n_d} & \to & \mathbb{N}^{(n_1 \cdots n_\mu) \times (n_{\mu+1} \cdots n_d)} \\ \operatorname{ind}(i_1, \ldots, i_d) & = & (i_{row}, i_{col}), \end{array}$$

where

$$\begin{split} i_{row} &= 1 + \sum_{s=1}^{\mu} (i_s - 1) \prod_{t=1}^{s-1} n_t, \\ i_{col} &= 1 + \sum_{s=\mu+1}^{d} (i_s - 1) \prod_{t=\mu+1}^{s-1} n_t \end{split}$$

GTT decomposition and matrix factorizations

We can compute the compression of the tensor by computing the SVD of the unfoldings.

Lemma (Oseledets 2011)

The **TT** rank of a tensor
$$\mathcal{T}$$
 is given by
 $\operatorname{tt-rank}(\mathcal{T}) = (\operatorname{rank}(\mathcal{T}^{<1>}), \dots, \operatorname{rank}(\mathcal{T}^{})).$

Input: Tensor
$$\mathcal{T}$$
, ranks k_1, \ldots, k_d)
Output: U_1, \ldots, U_d .
 $k_0 = k_d = 1$;
for $\mu = 1, \ldots, d - 1$ do
Reshape \mathcal{T} into $T^{<2>} \in \mathbb{C}^{k_{\mu-1}n_{\mu} \times (n_{\mu+1}...n_d)}$;
Compute rank- k_{μ} approximation $T^{<2>} \approx U\Sigma V^T$ (e.g. via SVD);
Reshape U into $U_{\mu} \in \mathbb{C}^{k_{\mu-1} \times n_{\mu} \times k_{\mu}}$;
Update \mathcal{T} via $T^{<2>} \leftarrow U^T X^{<2>} = \Sigma V^T$;

E The **proof** is

obtained by simply following the steps of the algorithm.

We can use *tolerances* instead of fixed ranks.

end

Set $U_d = \mathcal{T}$; Algorithm 1: TT-SVD $(\mathcal{T}, k_1, \dots, k_d)$

TT decomposition and matrix factorizations

And we can estimate the resulting error using the best approximation properties of the SVD.

Theorem (Oseledets 2011)

Let \mathcal{T}_{SVD} denote the tensor in TT decomposition obtained from TT-SVD. Then

$$\|\mathcal{T} - \mathcal{T}_{SVD}\| \leq \sqrt{\epsilon_1^2 + \dots + \epsilon_d^2}$$

where

$$\epsilon_{\mu}^{2} = \|T^{<\mu>} - U\Sigma V^{T}\|_{F}^{2} = \sigma_{k_{\mu}+1}^{2} + \sigma_{k_{\mu}+2}^{2} + \dots$$

- We can modify the algorithm to accommodate different compression algorithms than the SVD,
- We can also compute the approximation via sketching algorithms, and avoiding using all the entries of *T*.

If a vector of length $N = n_1 \times \ldots \times n_d$ is treated as a *d*-dimensional tensor with mode sizes n_k , and represented in TT-format, the matrices acting on it have the form

$$\mathcal{M}(i_1,\ldots,i_d,j_1,\ldots,j_d)=M_1(i_1,j_1)\ldots M(i_d,j_d), \qquad M_k(i_k,j_k)\in\mathbb{R}^{r_{k-1}\times r_k},$$

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*****: the first block indexes i_1, \ldots, i_d enumerate the rows, *****: the second block indexes j_1, \ldots, j_d enumerate the columns. Given \mathcal{M} in TT-format, and a vector \mathcal{X} in TT-format with cores X_k , and entries $X(j_1, \ldots, j_d)$ then the matrix-vector multiplication amounts to the following sum

$$\mathcal{Y}(i_1,\ldots,i_d)=\sum_{j_1,\ldots,j_d}\mathcal{M}(i_1,\ldots,i_d,j_1,\ldots,j_d)\mathcal{X}(j_1,\ldots,j_d)=Y_1(i_1)\ldots Y_d(i_d),$$

where $Y_k(i_k) = \sum_{j_k} M_k(i_k, j_k) \otimes X_k(j_k)$

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where $Y_k(i_k) = \sum_{j_k} M_k(i_k, j_k) \otimes X_k(j_k)$ The ranks of \mathcal{Y} are the product of the ranks of the matrix and of the vector! So we need to **recompress** after every matrix-vector product.

We can use the same routine as before to *represent* the two BVM matrices,

```
%% Time-dependent operator
kval = 5; % Grid power
m = 2<sup>kval</sup>; % Number of time
\hookrightarrow steps
k = 2;
[Alpha,Beta] = mab(k,m);
A = Alpha(:, 2:m+1);
B = Beta(:, 2:m+1);
t0 = 0:
tf = 1:
h = (tf-t0)/m:
tA = tt_matrix(full(A),1e-14);
tA = tt_reshape(tA,2*ones(kval,2));
tB = tt_eye(2, kval);
```

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- We build a tensor in which all the modes have size 2, this is usually called a Quantized-TT (QTT) formulation:

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If we look at the values of k and maximal tt-rank we find:

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```

We can act similarly also for the space operator.

%% Compression of the space part
tL = tt_matrix(L,1e-14);
tL = tt_reshape(tL,2*ones(kval+1,2));
tM = tt_eye(2,kval+1);
%% Final assembly
tMat = tkron(tA,tM)-h*tkron(tB,tL);

- We can act similarly also for the space operator.
- ▲ We could be way more clever in the representation of these matrices, these are diagonal times Toeplitz, and we could do something specialized, *e.g.*, (Kazeev, Khoromskij, and Tyrtyshnikov 2013).

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 - AMEn Use a *specialized solver* for linear systems in TT format (Dolgov and Savostyanov 2014).

Using AMEn (Dolgov and Savostyanov 2014) as

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2	64	128	9	2.231e-07	22
2	128	256	10	3.428e-07	26
2	256	512	14	5.925e-07	30
2	512	1024	22	3.957e-07	33
2	1024	2048	35	6.034e-07	37
2	2048	4096	47	6.968e-07	42

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3	128	256	11	2.153e-07	24
3	256	512	15	2.138e-07	28
3	512	1024	18	2.950e-07	32
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- Behavior is *similar* to the matrix-equation solver,
- **We** could play around with different **settings** and **options** of the AMEn solver.
- Studying the right combination of parameters, representation, setups is still an open problem for the BVM all-at-once approaches.

Conclusion and summary

- 🛇 We have seen how to work with linear multistep methods in boundary value form,
- We have discussed some structured preconditioning strategy for the resulting linear systems,
- We have introduced the machinery for working with tensor equations in the Tensor Train format.
- There are many open problems and possibilities to do better here.

Next up

- 📋 Fractional Laplacians,
- 📋 Rational approximations and matrix functions,
- B A couple of applications to complex network theory.
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