

# Solving Linear Systems of the Form $(A + \gamma UU^T)x = b$

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Iterative Methods for Large-Scale Saddle-Point Problems  
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# Palazzo della Carovana (1567), seat of the Scuola Normale



## Basic facts about the Scuola Normale

- Established by Napoleon in 1813 as a branch of the ENS in Paris
- Public institution (independent from U. of Pisa)
- Three Faculties: Sciences, Letters and Philosophy, Political and Social Sciences (approximately 65 permanent faculty members)
- Most selective educational institution in Italy (5% admission rate)
- Student body is nearly 600 students (half of them PhD students), all on full scholarships
- Some notable alumni: Giosuè Carducci (Nobel prize 1906), Enrico Fermi (Nobel prize 1938), Carlo Rubbia (Nobel prize 1984), Alessio Figalli (Fields medal 2018)
- Two of Italy's presidents and several prime ministers were SNS alumni

# Our conference room, the "Sala Azzurra"



- 1 Motivation
- 2 The proposed preconditioner
- 3 Eigenvalue bounds
- 4 Numerical experiments
- 5 Conclusions and future work

Joint work with Chiara Faccio (SNS).

Please note that this is [work in progress](#)!

# Outline

- 1 Motivation
- 2 The proposed preconditioner
- 3 Eigenvalue bounds
- 4 Numerical experiments
- 5 Conclusions and future work

# The problem

We are interested in finding efficient solvers for large systems of the form

$$(A + \gamma UU^T)x = b, \quad (1)$$

where  $A \in \mathbb{R}^{n \times n}$ ,  $U \in \mathbb{R}^{n \times k}$ ,  $\gamma > 0$  and  $b \in \mathbb{R}^n$ . Here  $1 \ll k \ll n$ .

We target problems with the following characteristics:

- $A$  is possibly singular, but  $A + \gamma UU^T$  is nonsingular for  $\gamma > 0$ .
- $A$  has one or more desirable property (sparsity, structure, etc.) which is lost if we form  $A + \gamma UU^T$  explicitly.
- Mat-vec products with  $A + \gamma UU^T$  can be computed efficiently.
- $k$  may not be "small", but  $k \times k$  systems can be solved accurately.
- Problem (1) must be solved repeatedly within a given application. Often, either  $A$  or  $U$  remains constant.

## The problem (cont.)

Problems of the form (1) which such characteristics arise frequently in scientific computing.

Examples include:

- Augmented Lagrangian methods for saddle point problems;
- Solution of KKT systems in constrained optimization;
- Solution of sparse-dense least squares problems;
- Certain types of integro-differential equations;
- Solution of PDEs describing almost incompressible materials;
- Numerical solution of PDEs with nonlocal BC's;
- ...



## Example 1: Augmented Lagrangian methods

Consider the saddle point problem

$$\mathcal{A} \mathbf{x} = \begin{bmatrix} A & B^T \\ B & 0 \end{bmatrix} \begin{bmatrix} u \\ p \end{bmatrix} = \begin{bmatrix} f \\ g \end{bmatrix} = \mathbf{f}.$$

Such systems arise frequently from the [finite element discretization](#) of systems of PDEs, such as for example the Stokes equations, the Oseen problem (obtained from the steady Navier-Stokes equations via Picard linearization), or the coupled Stokes-Darcy system.

A powerful approach to solve such systems is the one based on the [augmented Lagrangian](#).

M. Benzi and M. Olshanskii, *An augmented Lagrangian-based approach to the Oseen problem*, SIAM J. Sci. Comput., 28 (2006), pp. 2095–2113.

P. E. Farrell, L. Mitchell, and F. Wechsung, *An augmented Lagrangian preconditioner for the 3D stationary incompressible Navier-Stokes equations at high Reynolds numbers*, SIAM J. Sci. Comput., 41 (2019), pp. A3075–A3096.

## Example 1: Augmented Lagrangian methods (cont.)

The idea is to replace the original saddle point problem with an equivalent one of the form:

$$\mathcal{A}_\gamma \mathbf{x} = \begin{bmatrix} A + \gamma B^T W^{-1} B & B^T \\ B & 0 \end{bmatrix} \begin{bmatrix} u \\ p \end{bmatrix} = \begin{bmatrix} \hat{f} \\ g \end{bmatrix} = \hat{\mathbf{f}},$$

where  $\hat{f} = f + \gamma B^T W^{-1} g$ . Here  $W$  is usually diagonal and positive definite. In the finite element setting,  $W$  is often the diagonal of the (pressure) mass matrix.

This new, augmented system is then solved by a Krylov subspace method like (F)GMRES with preconditioner

$$\mathcal{P}_\gamma = \begin{bmatrix} A + \gamma B^T W^{-1} B & B^T \\ 0 & -\gamma^{-1} W \end{bmatrix}.$$

In practice, the preconditioner is applied **inexactly**.

## Example 1: Augmented Lagrangian methods (cont.)

The convergence of the preconditioned iteration is usually **very fast** and independent of parameters like the mesh size and viscosity, especially in the “large  $\gamma$ ” limit.

However, at each iteration of the Krylov subspace method a linear system with coefficient matrix  $A + \gamma B^T W^{-1} B$  must be solved (inexactly).

This linear system is of the form (1) with  $U = B^T W^{-1/2}$ . Here  $A$  is **sparse**, often **block diagonal**, and **positive definite** (or  $A + A^T$  is).

Forming  $A + \gamma B^T W^{-1} B$  explicitly would lead to **loss of sparsity and structure**. This system can be quite ill-conditioned (esp. for large  $\gamma$ ) and its solution is the **main challenge** associated with the augmented Lagrangian approach.

It is therefore necessary to develop efficient iterative methods for it.

Ideally, we would like such solvers to be robust with respect to  $\gamma > 0$ .

## Example 2: KKT systems in constrained optimization

The solution of (smooth) constrained minimization problems by interior point (IP) methods leads to sequences of linear systems of the form

$$\mathbf{A} \mathbf{x} = \begin{bmatrix} H & -C^T & 0 \\ C & 0 & -I \\ 0 & Z & \Lambda \end{bmatrix} \begin{bmatrix} \delta x \\ \delta \lambda \\ \delta z \end{bmatrix} = \begin{bmatrix} -r_1 \\ -r_2 \\ -r_3 \end{bmatrix} = \mathbf{f}.$$

Here  $H = H^T$  is the Hessian of the objective function at the current point  $x_k$ ,  $C$  is the Jacobian of the constraints at the same point, and  $Z$  and  $\Lambda$  are diagonal, positive definite matrices associated with the current values of the Lagrange multipliers  $\lambda_k$  and slack variables  $z_k$ , respectively.

The variable  $\delta z$  can easily be obtained using the last equation:

$$\delta z = -\Lambda^{-1}(r_3 + Z\delta\lambda)$$

and substituted into the second (block) equation.

## Example 2: KKT systems in constrained optimization (cont.)

This yields the reduced system

$$\begin{bmatrix} H & -C^T \\ C & \Lambda^{-1}Z \end{bmatrix} \begin{bmatrix} \delta x \\ \delta \lambda \end{bmatrix} = \begin{bmatrix} -r_1 \\ -r_2 - \Lambda^{-1}r_3 \end{bmatrix}.$$

Eliminating  $\delta \lambda$  leads to the **fully reduced** (Schur complement) system

$$(H + C^T Z^{-1} \Lambda C) \delta x = -r_1 - C^T Z^{-1} (r_3 + \Lambda r_2) =: b.$$

After solving for  $\delta x$ , the other unknowns  $\delta \lambda$  and  $\delta z$  are readily obtained.

This system is of the form (1) with  $A = H$ ,  $U = C^T (Z^{-1} \Lambda)^{1/2}$  and  $\gamma = 1$ .

The Hessian is usually **positive semidefinite, sparse and possibly structured**. Again, forming  $H + C^T Z^{-1} \Lambda C$  explicitly is generally undesirable. Instead, we propose to solve the fully reduced system iteratively, using a suitable (algebraic) preconditioner.

## Example 3: sparse-dense LS problems

Consider a large LS problem of the form

$$\|Bx - c\|_2 = \min,$$

where  $B \in \mathbb{R}^{m \times n}$  and  $c \in \mathbb{R}^m$ . Assume that  $B$  has the following structure:

$$B = \begin{bmatrix} B_1 \\ B_2 \end{bmatrix}, \quad B_1 \in \mathbb{R}^{(m-k) \times n}, \quad B_2 \in \mathbb{R}^{k \times n},$$

where  $B_1$  is **sparse** and  $B_2$  is **dense**. Then the LS problem is equivalent to the  $n \times n$  system of normal equations:

$$(B_1^T B_1 + B_2^T B_2)x = B^T c,$$

which is of the form (1) with  $A = B_1^T B_1$ ,  $U = B_2^T$ ,  $\gamma = 1$  and  $b = B^T c$ .

Once again, we would like to solve this system by an iterative method. The main challenge is again constructing an effective preconditioner.

# Outline

- 1 Motivation
- 2 The proposed preconditioner**
- 3 Eigenvalue bounds
- 4 Numerical experiments
- 5 Conclusions and future work

# The preconditioner

Consider again the linear system (1):  $(A + \gamma UU^T)x = b$ .

As we have seen, in applications the matrix  $A$  (or  $A + A^T$ ) is usually at least positive semidefinite, and we will make this assumption.

Also, although this is not strictly necessary, we will assume that

$$\text{Ker}(A) \cap \text{Ker}(U^T) = \{0\},$$

so that  $A + \gamma UU^T$  is nonsingular (and positive definite) for all  $\gamma > 0$ .

When  $A$  is **nonsingular**, we could use the Sherman-Morrison-Woodbury (SMW) formula to solve (1), but this is only applicable to problems of **moderate size**. Recall that SMW states that

$$(A + \gamma UU^T)^{-1} = A^{-1} - \gamma A^{-1}U(I_k + \gamma U^T A^{-1}U)^{-1}U^T A^{-1}.$$

Another possibility would be to build preconditioners based on the SMW formula, where the action of  $A^{-1}$  is replaced by some approximation, but our attempts were unsuccessful. Also, often  $A$  is **singular**.



## The preconditioner (cont.)

When  $k$  is small (say,  $k = 10$  or less) then any good preconditioner for  $A$  (or  $A + \alpha I_n$ ,  $\alpha > 0$ , if  $A$  is singular) tends to give good results. In fact, using CG preconditioned with  $A^{-1}$  yields convergence in at most  $k + 1$  steps. However, if  $k$  is in the hundreds (or larger), this approach is not appealing.

Hence, we need to take into account both  $A$  and  $\gamma U U^T$  when building the preconditioner. We do this by forming a suitable **product preconditioner**, as follows.

Let  $\alpha > 0$  and consider the two splittings

$$A + \gamma U U^T = (A + \alpha I_n) - (\alpha I_n - \gamma U U^T)$$

and

$$A + \gamma U U^T = (\alpha I_n + \gamma U U^T) - (\alpha I_n - A).$$

Note that both  $A + \alpha I_n$  and  $\alpha I_n + \gamma U U^T$  are invertible.

## The preconditioner (cont.)

Let  $x^{(0)} \in \mathbb{R}^n$  and consider the [alternating iteration](#)

$$(A + \alpha I_n)x^{(k+1/2)} = (\alpha I_n - \gamma U U^T)x^{(k)} + b,$$

$$(\alpha I_n + \gamma U U^T)x^{(k+1)} = (\alpha I_n - A)x^{(k+1/2)} + b,$$

with  $k = 0, 1, \dots$ . This alternating scheme is analogous to that of other well-known iterative methods like ADI, HSS, etc.

**Theorem 1:** Assume  $A + A^T$  is positive definite. Then the sequence  $\{x^{(k)}\}$  converges, as  $k \rightarrow \infty$ , to the unique solution of equation (1), for any choice of  $x^{(0)}$  and for all  $\alpha > 0$ .

To turn this into a practical method, we will use it as a [preconditioner](#) for a Krylov-type method rather than as a stationary iterative scheme. This will also allow [inexact solves](#).

## The preconditioner (cont.)

To derive the preconditioner we eliminate  $x^{(k+1/2)}$  and write the iterative scheme as the fixed-point iteration

$$x^{(k+1)} = T_\alpha x^{(k)} + c = (I_n - P_\alpha^{-1} A_\gamma) x^{(k)} + P_\alpha^{-1} b,$$

where we have set  $A_\gamma = A + \gamma U U^T$ . An easy calculation reveals that the preconditioner  $P_\alpha$  is given, in factored form, by

$$P_\alpha = \frac{1}{2\alpha} (A + \alpha I_n) (\alpha I_n + \gamma U U^T).$$

The scalar factor  $\frac{1}{2\alpha}$  is immaterial for preconditioning, and can be ignored.

Applying this preconditioner requires **two solves** involving  $A + \alpha I_n$  and  $\alpha I_n + \gamma U U^T$  at each Krylov iteration.

Generally speaking, each of these should be considerably **simpler** than solving systems involving the matrix  $A_\gamma = A + \gamma U U^T$ .

## The preconditioner (cont.)

Consider first solves involving  $A + \alpha I_n$ . If  $A$  is sparse, and/or structured (e.g., block diagonal, Toeplitz, etc.) then so is  $A + \alpha I_n$ .

Exact solves with  $A + \alpha I_n$  can be replaced, if necessary, with inexact solves using either a good preconditioner for  $A + \alpha I_n$  or a few steps of an inner iteration (PCG, AMG, or other).

Note the usual trade-off: larger values of  $\alpha$  make solves with  $A + \alpha I_n$  easier, but may degrade the performance of the preconditioner  $P_\alpha$ .

Numerical experiments suggest that the solution of linear systems involving  $\alpha I_n + \gamma U U^T$  is **more critical**. Note that this matrix is SPD for all  $\alpha > 0$ , but ill-conditioned for small  $\alpha$  (or very large  $\gamma$ ).

## The preconditioner (cont.)

The Sherman-Morrison-Woodbury formula yields

$$(\alpha I_n + \gamma U U^T)^{-1} = \alpha^{-1} I_n - \alpha^{-1} \gamma U (\alpha I_k + \gamma U^T U)^{-1} U^T.$$

The main cost is the solution at each step of a  $k \times k$  linear system with matrix  $\alpha I_k + \gamma U^T U$ , which can be performed by Cholesky factorization (computed once and for all at the outset) or possibly by a suitable inner PCG iteration or maybe an (algebraic) MG method.

Note that for incompressible flow problems,  $\alpha I_k + \gamma U^T U$  is essentially a (shifted) discrete pressure Laplacian.

In the numerical solution of the Navier–Stokes equations using (say) Picard iteration, this matrix remains constant, whereas the matrix  $A$  changes.

Hence, the cost of a Cholesky factorization can be **amortized** over many nonlinear (or time) steps.

## Two variants

Building on the main idea, some variants of the preconditioner can be envisioned.

If  $A$  happens to be nonsingular and linear systems with  $A$  are not too difficult to solve (e.g., well-conditioned), then it may not be necessary to shift  $A$ , leading to a preconditioner of the form

$$P_{\alpha,0} = A(\alpha I_n + \gamma U U^T).$$

Note that Theorem 1, however, is no longer applicable.

When  $A$  is symmetric positive semidefinite, we'd like the preconditioner to be SPD so that it can be used with the CG method. In this case we can consider a symmetrized version of the preconditioner, for example

$$P_{\alpha}^S = L(\alpha I + \gamma U U^T)L^T$$

where  $L$  is the Cholesky (or incomplete Cholesky) factor of  $A + \alpha I$  (or of  $A$  itself if  $A$  is SPD and not very ill-conditioned).

## Possible variants (cont.)

In some cases (but not always) the performance of the method improves if  $A_\gamma$  is diagonally scaled so that it has unit diagonal prior to forming the preconditioner.

Note that the matrix

$$D_\gamma = \text{diag}(A + \gamma U U^T)$$

can be easily computed:

$$(D_\gamma)_{ii} = a_{ii} + \gamma \|u_i^T\|_2^2,$$

where  $u_i^T$  is the  $i$ th row of  $U$ .

It is easy to see that applying the preconditioner to the diagonally scaled matrix  $D_\gamma^{-1/2} A_\gamma D_\gamma^{-1/2}$  is equivalent to using the modified preconditioner

$$\hat{P} = (A + \alpha D_\gamma)(\alpha D_\gamma + \gamma U U^T)$$

on the original matrix. Other choices of  $D_\gamma$  are also possible.

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## Bounds on the eigenvalues

Let  $A_\gamma := A + \gamma U U^T$  and  $P_\alpha := \frac{1}{2\alpha}(A + \alpha I)(\alpha I + \gamma U U^T)$ .

WLOG we can assume that  $\|A\|_2 = 1$  and  $\|U\|_2 = 1$ . We also assume that  $A_\gamma$  is nonsingular (that is,  $\text{Ker}(A) \cap \text{Ker}(U^T) = \{0\}$ ).

**Theorem 2.** Let  $A + A^T$  be positive semidefinite. If  $(\lambda, x)$  is an eigenpair of the preconditioned matrix  $P_\alpha^{-1} A_\gamma$ , with  $\|x\|_2 = 1$ , then

$$\mu < \text{Re}(\lambda) < 2, \quad |\text{Im}(\lambda)| < 1 \quad (2)$$

where

$$\mu = \frac{\alpha \lambda_{\min}(A + A^T)}{(1 + \alpha)(\alpha + \gamma)}.$$

If  $(\lambda, x)$  is an eigenpair with  $x \in \text{Ker}(U^T)$ , then

$$\lambda = \frac{2x^* A x}{x^* A x + \alpha}$$

(independent of  $\gamma$ ).

## Bounds on the eigenvalues (cont.)

Some comments on this result are in order:

- We see from (2) that the lower bound is uninformative if  $A + A^T$  is singular ( $\mu = 0$ ).
- The lower bound (if  $\neq 0$ ) is maximized for  $\alpha = \sqrt{\gamma}$ .
- Choosing  $\alpha = \sqrt{\gamma}$  to maximize the lower bound may not be optimal.
- The lower bound approaches 0 if  $\gamma \rightarrow \infty$ , indicating that the case of large  $\gamma$  may be challenging.
- The result assumes the preconditioner is applied **exactly** (often not true in practice).
- Eigenvalues alone may not be descriptive of GMRES convergence.

## Bounds on the eigenvalues (cont.)

$\gamma$	$\alpha$	$\max(\operatorname{Re}(\lambda))$	$\min(\operatorname{Re}(\lambda))$	lower bound
0.1	0.1	1.818e+00	1.700e-02	5.709e-04
	<b>0.3162</b>	1.519e+00	5.409e-03	7.250e-04
	5.0	3.333e-01	3.430e-04	2.052e-04
1.0	0.5	1.333e+00	6.590e-03	2.791e-04
	<b>1.0</b>	1.000e+00	3.300e-03	3.140e-04
	5.0	4.683e-01	6.609e-04	1.744e-04
50.0	1.0	1.532e+00	3.323e-03	1.231e-05
	<b>7.0711</b>	1.658e+00	4.707e-04	1.928e-05
	10.0	1.606e+00	3.328e-04	1.903e-05

**Table:** Stokes problem from IFISS,  $64 \times 64$  mesh and Q2-Q1 discretization.  $A$  and  $U$  normalized so that  $\|A\|_2 = 1 = \|U\|_2$ . In boldface the value  $\alpha = \sqrt{\gamma}$ .

## Bounds on the eigenvalues (cont.)

$\alpha$	$\max(\operatorname{Re}(\lambda))$	$\min(\operatorname{Re}(\lambda))$	lower bound
0.001	1.998e+00	6.508e-03	7.343e-04
0.01	1.980e+00	6.321e-02	7.213e-03
0.1	1.818e+00	4.834e-01	6.081e-02
0.5	1.333e+00	8.484e-01	1.635e-01
<b>1.0</b>	1.000e+00	5.384e-01	1.839e-01
5.0	4.335e-01	1.372e-01	1.022e-01
10.0	2.457e-01	7.106e-02	6.081e-02
20.0	1.313e-01	3.617e-02	3.337e-02
50.0	5.470e-02	1.463e-02	1.414e-02

**Table:** Problem `mosarqp1` from Maros and Mészáros collection (Schur complement of KKT system form constrained optimization),  $\gamma = 1$ .  $A$  and  $U$  normalized so that  $\|A\|_2 = 1 = \|U\|_2$ .

## Bounds on the eigenvalues (cont.)

$\alpha$	$\max(\operatorname{Re}(\lambda))$	$\min(\operatorname{Re}(\lambda))$	lower bound
0.001	1.998e+00	4.247e-03	1.597e-03
0.01	1.980e+00	4.167e-02	1.568e-02
0.1	1.818e+00	3.477e-01	1.322e-01
0.5	1.333e+00	9.087e-01	3.556e-01
<b>1.0</b>	1.000e+00	8.889e-01	4.000e-01
5.0	5.383e-01	2.759e-01	2.222e-01
10.0	3.180e-01	1.481e-01	1.322e-01
20.0	1.738e-01	7.692e-02	7.256e-02
50.0	7.350e-02	3.150e-02	3.076e-02

**Table:** Problem `1p_fit2p` from SuiteSparse Matrix Collection (sparse-dense least-squares problem),  $\gamma = 1$ .  $A$  and  $U$  normalized so that  $\|A\|_2 = 1 = \|U\|_2$ .

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# Numerical experiments with matrices from Stokes and Oseen problems (leaky-lid driven cavity)

We tested **inexact** variants of the proposed preconditioner

$$P_\alpha = \frac{1}{2\alpha}(A + \alpha I)(\alpha I + \gamma U U^T)$$

on a number of linear systems of the form

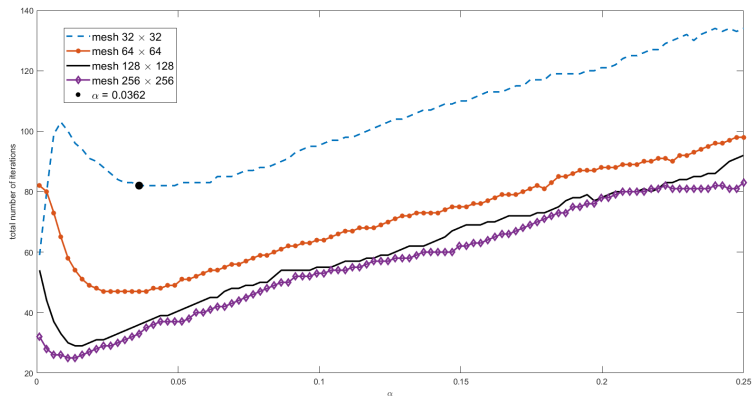
$$(A + \gamma B^T W^{-1} B)x = b$$

associated with 2D steady Stokes and Oseen problems, varying  $\gamma$ ,  $\alpha$ , the mesh size  $h$ , the viscosity  $\nu$  and the type of discretization used. Note that  $U = B^T W^{-1/2}$  and that  $W$  is diagonal. Also,  $A$  is **block diagonal**.

For efficiency, we replace the factor  $(A + \alpha I)$  with its no-fill Cholesky or ILU factorization, denoted by  $M_\alpha$ . The factor  $(\alpha I + \gamma B^T W^{-1} B)$  is inverted exactly via the SMW formula and the Cholesky factorization of the  $k \times k$  matrix  $\alpha I_k + \gamma W^{-1/2} B B^T W^{-1/2}$ .

We also describe a simple but effective heuristic for the choice of  $\alpha$ .

# Numerical experiments with steady 2D Oseen problem



**Figure:** Number of PGMRES iterations versus  $\alpha$  for the 2D Oseen problem (leaky-lid driven cavity from IFISS) with  $\nu = 0.01$ ,  $\gamma = 100$ , Q2-Q1 finite element discretization and different mesh sizes. GMRES restart  $m = 20$ , convergence residual tolerance =  $1e-06$ . Diagonal scaling is applied.

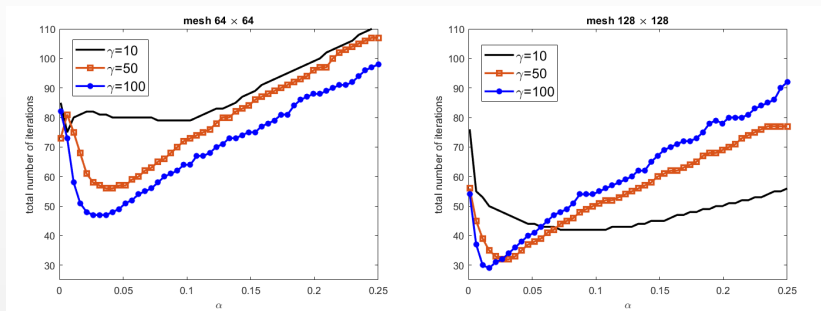


## A simple heuristic for the choice of $\alpha$

mesh	$\alpha^*$	iterations with $\alpha^*$	optimal $\alpha$	iterations with optimal $\alpha$
$64 \times 64$	0.0256	47	0.0236	47
$128 \times 128$	0.0181	30	0.0136	29
$256 \times 256$	0.0128	25	0.0111	25

**Table:** PGMRES iteration counts for 2D Oseen problem with  $\nu = 0.01$ ,  $\gamma = 100$ , Q2-Q1 finite element discretization. For  $32 \times 32$  mesh we find  $\alpha^* = 0.0362$ , then for the  $2^{5+k} \times 2^{5+k}$  mesh we set  $\alpha^* = \frac{0.0362}{2^{k/2}}$ . That is, we divide  $\alpha^*$  by  $\sqrt{2}$  each time  $h$  is halved.

# Numerical experiments for steady 2D Oseen problem



**Figure:** Number of iterations versus  $\alpha$  for the 2D Oseen problem from IFISS with  $\nu = 0.01$ , Q2-Q1 finite element discretization on  $64 \times 64$  mesh (LEFT) and on  $128 \times 128$  mesh (RIGHT) for different values of  $\gamma$ .

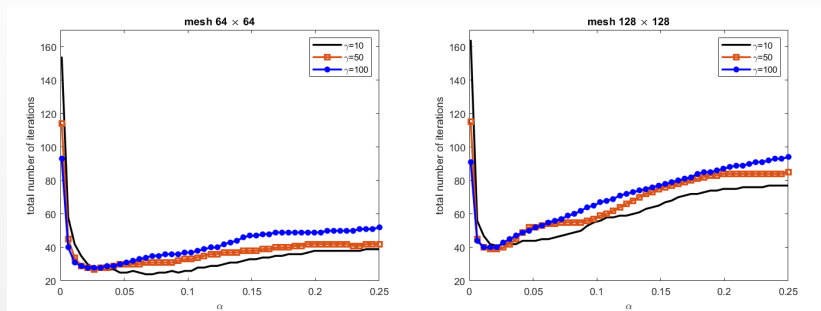
# Numerical experiments: timings for Oseen problem

mesh	M-Time	P-Time	$M_\alpha$		$P_\alpha$	
			Sol-Time	Its	Sol-Time	Its
$32 \times 32$	8.91e-04	1.55e-03	4.21e-01	1476	3.88e-02	82
$64 \times 64$	1.79e-03	4.13e-03	1.82e+00	1408	9.11e-02	47
$128 \times 128$	6.58e-03	2.02e-02	9.50e+00	1516	3.26e-01	30
$256 \times 256$	2.52e-02	1.16e-01	4.90e+01	1285	1.61e+00	25

**Table:** Linear system from Oseen problem with  $\gamma = 100$ ,  $\nu = 0.01$ , Q2-Q1 finite element discretization. Diagonal scaling is used.  $M_\alpha$  is ILU(0) of  $A + \alpha I$ ,  $P_\alpha = M_\alpha(\alpha I_k + \gamma U^T U)$  with Cholesky factorization of  $k \times k$  matrix in SMW formula. M-Time and P-Time are the preconditioner construction times. Sol-Time is the time for the preconditioned iteration to converge.

**Note:** In this application, the Cholesky factorization needed for inverting  $\alpha I + \gamma U^T U$  can be reused many times since  $U$  is constant.

# Numerical experiments for steady 2D Stokes problem



**Figure:** Number of iterations versus  $\alpha$  for the 2D Stokes problem from IFISS with Q2-Q1 finite element discretization on  $64 \times 64$  mesh (LEFT) and on  $128 \times 128$  mesh (RIGHT) for different values of  $\gamma$ .

## Numerical experiments on reduced KKT systems, I

Problem primal4 (Maros and Mészáros collection),  $n = 1489$ ,  $k = 75$ .  
The reduced system matrix  $H + C^T(Z^{-1}\Lambda)C$  has condition number  $3.41 \times 10^5$ . **Note:**  $H$  is **singular**.

$\alpha$	$M_\alpha$	$P_\alpha$	no prec.
0.001	2000*	13	2000*
0.01	1458	13	
0.1	1379	11	
0.5	1358	9	
1.0	2000*	2	
10.0	2000*	15	
20.0	2000*	18	

**Table:** Total number of iterations for primal4.  $M_\alpha$  is IC(0) of  $H + \alpha I$ . No diagonal scaling is applied.

## Numerical experiments on reduced KKT systems, II

Problem `mosarqp1` problem (Maros and Mészáros collection),  $n = 2500$ ,  $k = 700$ . The reduced system matrix  $H + C^T(Z^{-1}\Lambda)C$  has condition number  $3.35 \times 10^4$ .

$\alpha$	$M_\alpha$	$P_\alpha$	no prec.
0.001	2000*	309	2000*
0.01	2000*	66	
0.1	2000*	20	
0.5	2000*	9	
1.0	2000*	6	
10.0	2000*	11	
20.0	2000*	13	

$\alpha$	$P_\alpha^S$	no prec. CG
0.001	293	246
0.01	125	
0.1	45	
0.5	20	
1.0	14	
10.0	15	
20.0	17	

**Table:** Total number of iterations for `mosarqp1`. No diagonal scaling is applied. (LEFT) PGMRES. (RIGHT) PCG. We consider the symmetrized version of the preconditioner:  $P_\alpha^S = L(\alpha I_n + \gamma U U^T) L^T$ , where  $L$  is the no-fill incomplete Cholesky factor of  $H + \alpha I_n$ .

# Numerical experiments on sparse-dense LS problems, I

Problem `1p_fit2p` (SuiteSparse Collection). Here  $B_1$  is  $13500 \times 3000$ ,  $B_2$  is  $25 \times 3000$  (hence  $n = 3000$ ,  $k = 25$ ),  $\kappa(B_1^T B_1 + B_2^T B_2) = 2.52 \times 10^9$ .

**Note:** the norm of  $B_2^T B_2$  is 7 orders of magnitude larger than that of  $B_1^T B_1$ .

$\alpha$	$M_\alpha$	$P_\alpha$	no prec.
0.001	109	10	174
0.01	109	10	
0.1	109	10	
0.5	110	10	
1.0	110	8	
10.0	139	7	
20.0	149	9	

**Table:** Total number of (P)GMRES iterations for `1p_fit2p`.  $M_\alpha$  is IC(0) approximation of  $B_1^T B_1 + \alpha I$ . No diagonal scaling is used.

We remark that the cost for  $P_\alpha$  is only slightly larger than for  $M_\alpha$ .

# Numerical experiments on sparse-dense LS problems, II

Problem stormg2-1000.

Here  $B_1$  is  $1377185 \times 528185$ ,  $B_2$  is  $121 \times 528185$  (hence we have  $n = 528185$ ,  $k = 121$ ).

$\alpha$	$M_\alpha$	$P_\alpha$	no prec.
0.001	2000*	2000*	2000*
0.01	2000*	334	
0.1	2000*	98	
0.5	2000*	43	
1.0	2000*	50	
5.0	2000*	89	
10.0	2000*	118	
20.0	2000*	176	

**Table:** Total number of (P)GMRES iterations for stormg2-1000.  $M_\alpha$  is as before. No diagonal scaling is used.



# Numerical experiments: timings for sparse-dense LS problem

$\alpha$	$M_\alpha$				$P_\alpha$	
	M-Time	P-Time	Sol-Time	Its	Sol-Time	Its
0.5	7.05e-02	7.09e-02	1.01e+02	2000*	2.28e+00	43
1.0	6.90e-02	6.93e-02	1.01e+02	2000*	2.61e+00	50
1.5	6.93e-02	6.96e-02	1.01e+02	2000*	3.10e+00	59

**Table:** stormg2\_1000 problem.  $M_\alpha$  is IC(0) of  $A = B_1^T B_1 + \alpha I$ . No diagonal scaling is needed. In PGMRES, we do not form the matrix  $A = B_1^T B_1$ , but we compute mat-vecs as  $B_1^T(B_1 x)$ . Here  $B_1$  is  $1377185 \times 528185$ ,  $B_2$  is  $121 \times 528185$  (hence we have  $n = 528185$ ,  $k = 121$ ).

# Numerical experiments on sparse-dense LS problems, III

Problem scfxm1-2r (SuiteSparse Collection):  $B_1$  is  $65886 \times 37980$ ,  $B_2$  is  $57 \times 37980$  (so  $n = 37980$ ,  $k = 27$ ),  $\kappa(B_1^T B_1 + B_2^T B_2) = 9.32 \times 10^6$ .

**Note:**  $A = B_1^T B_1$  is **singular**.

$\alpha$	$M_\alpha$	$P_\alpha$	no prec.	$\alpha$	$P_\alpha^S$	no prec. CG
0.001	1572	555	240	0.001	1331	184
0.01	693	91		0.01	415	
0.1	183	36		0.1	105	
0.5	154	39		0.5	58	
1.0	155	50		1.0	65	
10.0	213	141		10.0	109	

**Table:** Total number of iterations for scfxm1-2r problem. Diagonal scaling is applied.  $M_\alpha$  is IC(0) of  $B_1^T B_1 + \alpha I$ . Left: PGMRES. Right: PCG. We consider the symmetrized version of the preconditioner:  $P_\alpha^S = L(\alpha I_n + \gamma U U^T) L^T$ , where  $L$  is the no-fill incomplete Cholesky factor of  $B_1^T B_1 + \alpha I$ .

# Outline

- 1 Motivation
- 2 The proposed preconditioner
- 3 Eigenvalue bounds
- 4 Numerical experiments
- 5 Conclusions and future work

# Conclusions

- Introduced a new solver for a wide class of tough linear systems
- The proposed preconditioner seems to work well in practice
- Solves with  $A + \alpha I_n$  can be (very) inexact...
- ... but exact solves with  $\alpha I + \gamma U U^T$  seem to be necessary
- SMW formula  $\Rightarrow$  only a  $k \times k$  solve needed
- Often, much of this work can be reused
- For some PDE problems we found a simple heuristic for choosing  $\alpha$
- Some theory available for ideal case

- Further investigate the spectrum of  $P_\alpha^{-1}A_\gamma$ . Clustering?
- Try to find eigenvalue bounds for singular  $A$
- More work to be done on the choice of  $\alpha$
- Investigate other approaches to solving systems with  $\alpha I + \gamma UU^T$  (iterative?)
- Are there better ways to symmetrize the preconditioner for use with CG?
- Extension to matrices of the form  $A + \gamma UV^T$ , with  $V \neq U$ .